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Tuesday, June 4, 13

Thanks to Julien for inviting me.

It is great to be hear. I learned a lot, met a lot of great people, hope to meet the rest of you too.

Thanks to Magnus & other folks who had trouble with the network.

The CouchDB Implementation

CouchDB is a database that replicates.

«Think of CouchDB as `git` for your application data.»

- Jan Lehnardt, Berlin Buzzwords

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multiple locations / push / pull



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PouchDB

TouchDB



CORE FEATURES

API

CORE DATA STRUCTURES

FILE SYSTEM ACCESS

FILE SYSTEM

COUCH_HTTP*

COUCH_DOC COUCH_MR COUCH_REPL...

COUCH_BTREE

COUCH_FILE

FILE SYSTEM

Core Datastructures

- 2 x b+tree & data interleaved
- append only, mvcc
- full fsync control
- Can answer:
- Data for \$key
- What happened \$since
- Used for core data storage
- As well as indexes
- Everything else is built on top

Behold the b-tree



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A		C	Δ	H	4	DOC_G	H		ſ	Κ		:
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A	8	C	DOC_D	Ε	4	DOC_G	H		ſ	DOC_K	ľ	•••
---	---	---	-------	---	---	-------	---	--	---	-------	---	-----

DOC_D	DOCG	DOC_K
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by-sequence nappened since?"

1. DOC_G	2. DOC_D	3. DOC_K
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The CouchDB File Format

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HEADER

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HEADER DOC_A

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Tuesday, June 4, 13 Bulk add + Delete










Tuesday, June 4, 13 Bulk add + Delete



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Tuesday, June 4, 13 Bulk add + Delete







Tuesday, June 4, 13 Bulk add + Delete

Operational Consequences

- efficient on spinning disk, "tape"
- btree = wide, upper layers in disk cache
- backup with cp \$a \$b
- crash safety/recovery
- compaction hurts

Efficient With Storage

Cache Friendly

cp db.couch /mnt/backup



Core Features (using by-seq)

- Replication
- Indexing / Views / GeoCouch / Lucene / ES etc.
- /_changes
- Compaction

Operational Consequences

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- compaction hurts

Tuesday, June 4, 13 **Replication**

DATABASE A

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1 DATABASE A

Tuesday, June 4, 13





4	
3	
2	
1	
DATABASE A	







4	
3	3
2	2
1	1
DATABASE A	DATABASE B

4	4
3	3
2	2
1	1
DATABASE A	DATABASE B

5	
4	4
3	3
2	2
1	1
DATABASE A	DATABASE B

6	
5	
4	4
3	3
2	2
1	1
DATABASE A	DATABASE B

7	
6	
5	
4	4
3	3
2	2
1	
DATABASE A	DATABASE B









8	8
7	7
6	6
5	5
4	4
3	3
2	2
1	1
DATABASE A	DATABASE B

DATABASE A

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Indexing
1 DATABASE A

Tuesday, June 4, 13





4
3
2
1
DATABASE A







4	
3	3
2	2
1	1
DATABASE A	INDEX A

4	4
3	3
2	2
1	1
DATABASE A	INDEX A

5	
4	4
3	3
2	2
1	1
DATABASE A	INDEX A

6	
5	
4	4
3	3
2	2
1	1
DATABASE A	INDEX A

7	
6	
5	
4	4
3	3
2	2
1	1
DATABASE A	INDEX A







8	
7	7
6	6
5	5
4	4
3	3
2	2
1	1
DATABASE A	INDEX A

8	8
7	7
6	6
5	5
4	4
3	3
2	2
1	1
DATABASE A	INDEX A

DATABASE A

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1 DATABASE A

Tuesday, June 4, 13





4
3
2
1
DATABASE A

5
4
3
2
1
DATABASE A

6
5
4
3
2
1
DATABASE A

7
6
5
4
3
2
1
DATABASE A

8
7
6
5
4
3
2
1
DATABASE A

Tuesday, June 4, 13 Compaction

DATABASE A

Tuesday, June 4, 13

1. DOC_A

DATABASE A

Tuesday, June 4, 13

2. DOC_B 1. DOC_A DATABASE A

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Tuesday, June 4, 13 The Happy Path

DATABASE A

Tuesday, June 4, 13

DOC A [1-A] DATABASE A Tuesday, June 4, 13







DATABASE B

Tuesday, June 4, 13









DATABASE B

Tuesday, June 4, 13

DOC A [4-D,3-C,2-B,1-A]



DOC A [4-D,3-C,2-B,1-A]

DOC A [3-C,2-B,1-A]

DATABASE B

Tuesday, June 4, 13

Tuesday, June 4, 13 Conflicts!

DATABASE A

Tuesday, June 4, 13 Conflicts!

DOC A [1-A] DATABASE A Tuesday, June 4, 13







DATABASE B









DATABASE B

Tuesday, June 4, 13







DATABASE A

DOC A [4-D,3-C,2-B,1-A]

DOC A [3-C,2-B,1-A]

DATABASE B

Tuesday, June 4, 13

DOC A [[4-K,3-C,2-B,1-A], [4-D, 3-C,2-B,1-A]]

DOC A [4-K,3-C,2-B,1-A]

DOC A [3-C,2-B,1-A]

DOC A [2-B,1-A]

DOC A [1-A] DATABASE A DOC A [4-D,3-C,2-B,1-A]

DOC A [3-C,2-B,1-A]

DATABASE B

Tuesday, June 4, 13

DOC A [[4-K,4-D], 3-C,2-B,1-A]

DOC A [[4-K,3-C,2-B,1-A], [4-D, 3-C,2-B,1-A]]

DOC A [4-K,3-C,2-B,1-A]

DOC A [3-C,2-B,1-A]

DOC A [2-B,1-A]

DOC A [1-A] DATABASE A DOC A [4-D,3-C,2-B,1-A]

DOC A [3-C,2-B,1-A]

DATABASE B

Tuesday, June 4, 13

DOC A [4-K,4-D,3-C,2-B,1-A]

DOC A [[4-K,4-D], 3-C,2-B,1-A]

DOC A [[4-K,3-C,2-B,1-A], [4-D, 3-C,2-B,1-A]]

DOC A [4-K,3-C,2-B,1-A]

DOC A [3-C,2-B,1-A]

DOC A [2-B,1-A]

DOC A [1-A] DATABASE A DOC A [4-D,3-C,2-B,1-A]

DOC A [3-C,2-B,1-A]

DATABASE B

Tuesday, June 4, 13

DOC A [4-D,4-K,3-C,2-B,1-A]

DOC A [4-K,4-D,3-C,2-B,1-A]

DOC A [[4-K,4-D], 3-C,2-B,1-A]

DOC A [[4-K,3-C,2-B,1-A], [4-D, 3-C,2-B,1-A]]



DOC A [3-C,2-B,1-A]

DOC A [2-B,1-A]

DOC A [1-A] DATABASE A DOC A [4-D,3-C,2-B,1-A]

DOC A [3-C,2-B,1-A]

DATABASE B

Tuesday, June 4, 13

DOC A [5-L,4-D,4-K,3-C,2-B,1-A]

DOC A [4-D,4-K,3-C,2-B,1-A]

DOC A [4-K,4-D,3-C,2-B,1-A]

DOC A [[4-K,4-D], 3-C,2-B,1-A]

DOC A [[4-K,3-C,2-B,1-A], [4-D, 3-C,2-B,1-A]]



Tuesday, June 4, 13

- Small codebase
- Efficient in small teams
- Isolated processes
- Supervision tree
- Concurrency
- Portable runtime
- Hard to recruit for
- Steep ramp-on
- Rit of an operational black box (nine nines story)

Potential Improvements

- Smarter compactor
- Smarter file-storage
- Less custom HTTP handling
- More indexers



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